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09/828,188	04/09/2001	Kenneth Graham Paterson	1509-166	1046
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C/o Hewlett-Pa	ckard Company			
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Please find below and/or attached an Office communication concerning this application or proceeding.

7

		Application No.	Applicant(s)			
Office Action Summary		09/828,188	PATERSON, KENNETH GRAHAM			
		Examin r	Art Unit			
		Esaw T Abraham	2133			
Period fo	Th MAILING DATE of this communication app or Reply	pears on the cover she t with the c	orrespondence address			
A SH THE I - Exter after - If the - If NO - Failu - Any r	ORTENED STATUTORY PERIOD FOR REPL'MAILING DATE OF THIS COMMUNICATION. Insions of time may be available under the provisions of 37 CFR 1.13 SIX (6) MONTHS from the mailing date of this communication. In period for reply specified above is less than thirty (30) days, a reply operiod for reply is specified above, the maximum statutory period to reply within the set or extended period for reply will, by statute reply received by the Office later than three months after the mailing and patent term adjustment. See 37 CFR 1.704(b).	36(a). In no event, however, may a reply be tin y within the statutory minimum of thirty (30) day vill apply and will expire SIX (6) MONTHS from , cause the application to become ABANDONE	nely filed s will be considered timely. the mailing date of this communication. D (35 U.S.C. § 133).			
1)🖂	Responsive to communication(s) filed on 09 A	A <i>pril</i> 2001 .				
2a)[_	This action is FINAL . 2b)⊠ Th	is action is non-final.				
3)	Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under <i>Ex parte Quayle</i> , 1935 C.D. 11, 453 O.G. 213.					
·	ion of Claims					
•	Claim(s) 1 - 31 is/are pending in the application	•				
	4a) Of the above claim(s) is/are withdrawn from consideration.					
•	Claim(s) is/are allowed.					
	Claim(s) <u>1-31</u> is/are rejected.					
· · · · ·	Claim(s) is/are objected to.					
•	Claim(s) are subject to restriction and/o ion Papers	r election requirement.				
9)[The specification is objected to by the Examine	r.				
10) 🗌 .	The drawing(s) filed on is/are: a)□ accep	oted or b) objected to by the Exa	miner.			
_	Applicant may not request that any objection to the					
11) 🗀	The proposed drawing correction filed on		ved by the Examiner.			
40\□∶	If approved, corrected drawings are required in rep	•				
, —	The oath or declaration is objected to by the Ex	aminer.				
	under 35 U.S.C. §§ 119 and 120) (-l) (0			
•	Acknowledgment is made of a claim for foreigr	i priority under 35 U.S.C. § 119(a)-(a) or (i).			
a)(<i>` `</i>	s have been received				
	 1. Certified copies of the priority documents have been received. 2. Certified copies of the priority documents have been received in Application No 					
	3. Copies of the certified copies of the prior					
* 9	application from the International Bu See the attached detailed Office action for a list	reau (PCT Rule 17.2(a)).	•			
14) 🗀 A	Acknowledgment is made of a claim for domesti	c priority under 35 U.S.C. § 119(e) (to a provisional application).			
) \square The translation of the foreign language pro Acknowledgment is made of a claim for domest	• .				
Attachmen	t(s)					
2) Notic	e of References Cited (PTO-892) e of Draftsperson's Patent Drawing Review (PTO-948) mation Disclosure Statement(s) (PTO-1449) Paper No(s) <u>5</u>	5) Notice of Informal	/ (PTO-413) Paper No(s) Patent Application (PTO-152)			
S. Patent and T	rademark Office					

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DETAILED ACTION

1. Claims 1 to 31 are presented for examination.

Priority

Acknowledgment is made of applicant's claim for foreign priority under 35
 U.S.C. 119(a)-(d). The certified copy has been filed in parent Application No: 00303023.6
 (EPO) filed on 04/10/2000.

Information Disclosure Statement

3. The references listed in the information disclosure statement submitted on 04/09/01 have been considered. (See attached PTO-1449).

Claim Rejections - 35 USC § 112

The following is a quotation of the second paragraph of 35 U.S.C 112:

The specification shall conclude with one or more claims particularly pointing out and distinctly claiming the subject matter which the applicant regards as his invention.

- 4. Claims 1, 5, 9, 17, 21 and 25 recite the limitation "the roots of the polynomial used in the said polynomial remaindering" and "the roots of a generator polynomial" in claims 1, 5, 9, 17, 21 and 25. There is insufficient antecedent basis for this limitation in the claim.
- 5. Claims 1, 5, 9, 17, 21 and 25 are rejected under 35 U.S.C. 112, second paragraph, as being incomplete for omitting essential structural cooperative relationships of elements, such omission amounting to a gap between the necessary structural connections. See MPEP § 2172.01. The omitted structural cooperative relationships are: It is not clear where in the

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claimed method said "the roots of the polynomial used in said polynomial remaindering process are different to the roots of a generator polynomial of said Reed-Solomon error correcting code". It is not clear if the polynomial remaindering process and the Reed-Solomon error correcting code confined in the generator polynomial. The interconnection of such performance of checksum calculation with the generator polynomial can neither be visualized in the drawings nor can be clearly understood from the claimed language for proper examination purposes. The examiner would appreciate if the applicant would clarify this matter.

Claim Rejections - 35 USC § 103

The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negatived by the manner in which the invention was made.

The factual inquiries set forth in *Graham* v. *John Deere Co.*, 383 U.S. 1, 148 USPQ 459 (1966), that are applied for establishing a background for determining obviousness under 35 U.S.C. 103(a) are summarized as follows:

- 1. Determining the scope and contents of the prior art.
- 2. Ascertaining the differences between the prior art and the claims at issue.
- 3. Resolving the level of ordinary skill in the pertinent art.
- 4. Considering objective evidence present in the application indicating obviousness or nonobviousness.
- 6. Claims 1-31 are rejected under 35 U.S.C. 103(a) as being unpatentable over Tenengolts (U.S. PN: 4,782,490) in view of Lee et al. (U.S. PN: 5,872,799).

As per claims 1, 5, 16, 17 and 21, Tenengolts substantially teach or disclose a method or system for error detection and correction in which codewords are made up of data and two groups of check symbols. The first group of check symbols is generated by a correction verification code, which verifies error correction; and, the second group of check symbols is generated by an interleaved Reed-Solomon code with symbols from the Galois field GF (28), which serves for error correction and the correction verification code is cyclic with a generator polynomial over the GF (2⁸). The error correction system uses the first root of an error location polynomial to calculate the second root of the polynomial and the detection system, which employs a portion of the error correction system circuitry, uses a cyclic code with a generator polynomial is a root of a primitive polynomial over the GF (see abstract). Tenengolts do not explicitly teach a method of performing a checksum that includes a byte based polynomial remaindering. However, Lee et al. in an analogous art teach a check symbols that represent redundant information about the code word and used to provide error correction and detection capabilities and the check symbols are the coefficients of the remainder polynomial generated by dividing the order of polynomial by an order of "generator polynomial" over a Galois field (see col. 1, lines 46-65 and abstract). Further, Lee et al. teach a method of interleaving Reed-Solomon error correcting code word comprising the steps of receiving a plurality of data symbols; computing, in each interleave a code word including said data symbols and a plurality of check symbols, said check symbols being generated by a generator polynomial where alpha is a primitive element of GF (2^m) (see claim 1). Furthermore, Lee et al. teach a method of storing and retrieving data from a mass storage media packed in 8-bit bytes and a check symbol generator operates from the packed bytes (see col. 10, lines 40-55). Therefore, it would have

been obvious to a person having an ordinary skill in the art at the time the invention was made to implement the teachings of Tenengolts with the method of retrieving (reading out) data from a mass storage packed in 8-bit bytes (byte-based) and a check symbol generator operates from the packed byte as taught by Lee et al. **This modification** would have been obvious because a person having ordinary skill in the art would have been motivated to do so because it would be relatively high in operation to achieve a reduction in power consumption and an increase in speed of decoding operation.

As per claims 2, 6, 18 and 22, Tenengolts in view of Lee et al. teach all the subject matter claimed in claims 1, 5, 17 and 21 including Tenengolts teaches first group of check symbols generated by a correction verification code and second group of check symbols generated by Reed-Solomon code with symbols from the Galois field GF (2⁸), which serves for error correction and the correction verification code is cyclic with a generator polynomial over the GF (2⁸). Further, Lee et al. teach that receiving a plurality of data symbols and a plurality of check symbols, said check symbols being generated by a generator polynomial where alpha is a primitive element of GF (2^m). Tenengolts in view of Lee et al. do not explicitly teach a polynomial X²+X alpha²+alpha where alpha is the primitive element GF (2⁸). However, using polynomials having different degrees is common practice for most of polynomial generators and the implementation is up to the designers' choice depending on the requirement of the system to correct/detect errors. Therefore, it would have been obvious to a person having an ordinary skill in the art at the time the invention was made to employ a process of using a polynomial of different degrees for defining redundancy codes. This modification would have been obvious

because a person having ordinary skill in the art would have been motivated in order to permit flexibility of achieving higher coding gains and lower decoding complexities.

As per claims 3, 4, 7, 8, 19, 20, 23 and 24, Tenengolts in view of Lee et al. teach all the subject matter claimed in claims 1, 5, 17 and 21 including Lee et al. in figure 1 teach or disclose that in a storage system an error condition may exist which causes all symbols stored in the storage system to be zeroed and under that error condition, from a syndrome computation, the resulting syndrome coefficients become all zeroes, thereby *masking* the error condition (see col. 12, lines 12-20).

As per claims 9 and 25, Tenengolts substantially teach or disclose a method or system for error detection and correction in which codewords are made up of data and two groups of check symbols. The first group of check symbols is generated by a correction verification code, which verifies error correction; and, the second group of check symbols is generated by an interleaved Reed-Solomon code with symbols from the Galois field GF (2⁸), which serves for error correction and the correction verification code is cyclic with a generator polynomial over the GF (2⁸). The error correction system uses the first root of an error location polynomial to calculate the second root of the polynomial and the detection system, which employs a portion of the error correction system circuitry, uses a cyclic code with a generator polynomial is a root of a primitive polynomial over the GF (see abstract). Further, Tenengolts teach that a system for correcting errors in a codeword that includes data, a first group of check symbols and a second group of check symbols, at least a portion of the first group of check symbols having been generated by a code with a generator polynomial over the Galois field GF (2^m) the second group

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of check symbols having been generated from the data and the first group of check symbols as "data" by an interleaved Reed-Solomon code with a generator polynomial (see claim 25). Tenengolts do not explicitly teach performing a checksum includes a byte based polynomial remaindering. However, Lee et al. in an analogous art teach a check symbols that represent redundant information about the code word and used to provide error correction and detection capabilities and the check symbols are the coefficients of the remainder polynomial generated by dividing the order of polynomial by an order of "generator polynomial" over a Galois field (see col. 1, lines 46-65 and abstract). Further, Lee et al. teach a method of interleaving Reed-Solomon error correcting code word comprising the steps of receiving a plurality of data symbols; computing, in each interleave a code word including said data symbols and a plurality of check symbols, said check symbols being generated by a generator polynomial where alpha is a primitive element of GF (2^m) (see claim 1). Furthermore, Lee et al. teach a method of storing and retrieving data from a mass storage media packed in 8-bit bytes and a check symbol generator operates from the packed bytes (see col. 10, lines 40-55). Therefore, it would have been obvious to a person having an ordinary skill in the art at the time the invention was made to implement the teachings of Tenengolts with the method of retrieving (reading out) data from a mass storage packed in 8-bit bytes (byte-based) and a check symbol generator operates from the packed byte as taught by Lee et al. This modification would have been obvious because a person having ordinary skill in the art would have been motivated to do so because it would be relatively high in operation to achieve a reduction in power consumption and an increase in speed of decoding operation.

As per claims 10 and 26, Tenengolts in view of Lee et al. teach all the subject matter claimed in claims 9 and 25 including Tenengolts teaches first group of check symbols generated by a correction verification code and second group of check symbols generated by Reed-Solomon code with symbols from the Galois field GF (28), which serves for error correction and the correction verification code is cyclic with a generator polynomial over the GF (2⁸). Further, Lee et al. teach that receiving a plurality of data symbols and a plurality of check symbols, said check symbols being generated by a generator polynomial where alpha is a primitive element of GF (2^m). Tenengolts in view of Lee et al. do not explicitly teach a polynomial X²+X.alpha²+alpha where alpha is the primitive element GF (2⁸). However, using polynomials having different degrees is common practice for most of polynomial generators and the implementation is up to the designers' choice depending on the requirement of the system to correct/detect errors. Therefore, it would have been obvious to a person having an ordinary skill in the art at the time the invention was made to employ a process of using a polynomial of different degrees for defining redundancy codes. This modification would have been obvious because a person having ordinary skill in the art would have been motivated in order to permit flexibility of achieving higher coding gains and lower decoding complexities.

As per claims 11, 12, 27 and 28, Tenengolts in view of Lee et al. teach all the subject matter claimed in claims 9 and 25 including Lee et al. in figure 1 teach or disclose that in a storage system an error condition may exist which causes all symbols stored in the storage system to be zeroed and under that error condition, from a syndrome computation, the resulting syndrome coefficients become all zeroes, thereby *masking* the error condition (see col. 12, lines 12-20).

As per claims 13, 15, 29 and 31, Tenengolts in view of Lee et al. teach all the subject matter claimed in claims 9 and 25. Tenengolts substantially teach a system for correcting errors in a codeword that includes data, a first group of check symbols and a second group of check symbols, at least a portion of the first group of check symbols having been generated by a code with a generator polynomial over the Galois field GF (2^m) the second group of check symbols having been generated from the data and the first group of check symbols as "data" by an interleaved Reed-Solomon code with a generator polynomial (see claim 25). Tenengolts in view of Lee et al. do not explicitly teach check sum calculations operate mis-correct error 1 in 2¹⁶. However, the technique of operating performing a check to operate mis-correct errors of any proportion (such as 1 in 2¹⁶ or any value) is up to the designer's choice depending the systems' requirement. Therefore, it would have been obvious to one of ordinary skill in the art at the time the invention was made to design or implement the size of checksum for detecting mis-correct errors. This motivation would have been obvious to one ordinary skill in the art at the time the invention was made because one of ordinary skill in the art would have employed a process of detecting mis-correct errors in order to improve the information bit rate and the efficiency of the system.

As per claims 14 and 30, Tenengolts in view of Lee et al. teach all the subject matter claimed in claims 9 and 25. Tenengolts substantially teach that an error detection code employed for a high data transfer rate, a basic polynomial of an interleaved Reed-Solomon code employed having a minimal number of nonzero coefficients and a generator polynomial of the code have a form and be irreducible over the Galois field GF (2^m) then a shift register for the encoding, a combinatorial circuit for multiplication and an encoder results for the generalized Hamming

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code, which corrects single symbol or detects double symbol errors (see col. 15, lines 60-67 and

col. 16, lines 1-67).

Conclusion

7. The prior art made of record and not relied upon is considered pertinent to applicant's

disclosure.

US PN: 4,413,339

Riggle et al.

US PN: 5,822,337

Zook et al.

Any inquiry concerning this communication or earlier communication from the examiner 8.

should be directed to Esaw Abraham whose telephone number is (703) 305-7743. The examiner

can normally be reached on M-F 8-5.

If attempts to reach the examiner by telephone are successful, the examiner's supervisor,

Albert DeCady can be reached on (703) 305-9595. The fax phone numbers for the organization

where this application or proceeding is assigned are (703) 746-7239 for regular communications

and (703) 746-7238 for after final communications.

Any inquiry of a general nature or relating to the status of this application or proceeding

should be directed to the receptionist whose telephone number is (703) 305-3900.

Esai Ollaham

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epuy J. Lamarre

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Primary Examiner